UNIT-II LINEAR DATA STRUCTURES - STACKS, QUEVES

Stace ADT - Operations - Applications - Evaluating arithmetic expressions - Conversion of Infix to postfix expression-Queue ADT - Operations - Circular Queue - Priority Queuedequeue - applications of queues.

Introduction to Stack; Stacks are known as LIFO Clast in, first out) lists. It is an ordered list of the same type of elements. A is an ordered list of the same type of elements. A is an ordered list where all insertions and deletions stack is a linear list where all insertions and deletions stack is a linear list where all insertions and deletions are permitted only at one end of the list. When the are permitted only at one end of the list. When the elements are added to stack it grows at one end. elements are added to stack it grows at one end. Similarly, when the elements core deleted from a stack, it shrinks at the same end. it shrinks at the same end. it insertions and deletions are performed is only Hue, insertions and deletions are performed the top: one perition, namely, the end of the list, called the top: one perition, namely, the end of the list, called the top:

& The fundamental operations on a stack are Pinh, which is equivalent to an insert, and Pop, which deletes the most recently inserted element.

<Pop(S) Stack S Push(X,S) (Top(S) Stack S Push(X,S) Fig. Stack model: input to a stack is by Push; output is by Pop.

to The most recently inserted element can be examined prior to performing a Pop by use of the Top roubine. & The model depicted is Fig. signifies only that Pinches are ifp operations and Pops and Tops are output.

Top 2 Fig. Stack model: only the top element is accessible. 4 1 3 Fig. shows an abstract stack after several operations. * The Top is med to keep track of the index of the top most element. Array Implementation of Stacks;-In array implementation, if Top is -1, then it is an A stack is defined as a pointer to a structure. The structure contains the Top Of Stack and Capacity fields. Once the maximum size is known, the stack amay can be dynamically allocated. # ifndet -Stack-h Shout Stack Record; typedef shout Stack Record & Stack; int Is Empty (Stack S); int Is Full (Stack S); Stack Create Stack (int Max Elementi); void Dispose Stack (Stack S); Void Make Empty (Stack S); void Push (Element Type X, Stack S); Element Type Top (Stack S); void Pop (Stack S); P ElementType TopAndPop(Stack S); #endif 1* _Stack_h*/

14 Place is implementation file &/ A Stack implementation is a dynamically allocated away of Adofine Empty TOS(-1) # define Min Stack Size (5) Somet Stack Record int Capacity; int TopOf Stack; j; Element Type * Array; Stack Creation: Stack Create Stack (int Mox Elementh) Stack S; if (Max Elements < Min Stack Size) Error ("Stack size is too small"); 1 & Allocation memory \$ S= malloc (size of (shult Stack Record)); if (S== NULL) " Out of Space!!!!"); Fatal Error (" Out of Space!!!!"); 12 Allocation of Stack Amay */ S-7 Array = malloc (size of (Element Type) & Max Elements); if (S > Array == NULL) Fatal Error ("Out of space!!!"); 14 Initialize Topof Stack and Capacity of 5-7 Capacity = Max Elements; Marke Empty (S); return sj grater.

Dispose the stack: 1.15 Void Dispose Stack (Stack S) 2 if (S! = NOLL) . . 2 free (S->Array); 3 3 free (S); This routine first frees the stack away and then the stack structure. JA Empty: int Is Empty (Stack S) I return S-> Top Of Stack == Emply TOS; Make Emply Void Make Empty (Stack S) ² S -> Top Of Stack = Empty TOS; void Push (Element Type X, Stack S) Push: if (In Fall (S)) Error ("Full stack"); S -> Array E++S -> Top Of Stack] = X; else Retain Top of Stack: Element Type Top (Stack S) if (! Is Empty (s)) return S-7 Array[S-1 TopOfStack];

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Error ("Empty Staci"); return o; 1t Return value med to avoid warning */ Pop: void pop(Stack S) if (In Empty (S)) Error ("Empty Stoce"); else S-7 Top Of Stace --; 2 Routine to give tops element and pop a stack: Element Type TepAnd Pop (Stack S) if (! In Froppy (S)) return S-> Array [S-> Top Of Stack --]; Error ("Empty Stace"); 2 return og In computer's memory, stacks can be represented as a linear It, which is used to store the address of the topmart element of the stark. att is the position where the element will be added to or of the stack. at Another variable MAX is used to store the maximum no. of elements that the stack can hold. deleted from. of If TOP = NULL, then it indicates the stack is empty and if TOP = MAX-1, then it indicates the stack is full. ABCDEFGHT 012345678 9 TOP=9

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Operations on Stack:

A stack supports three banic operation: push, pop and prek. pinh => adds an element to the top of the stack. p.p =7 removes the element from " " " peer of seturn the value of the tepmost element of stack. Pop: Push: Als: A18: D. If TOP=NOLL 1). If TOP=MAX-1 Error " Under flow" Error "Overflow" gito step 4 goto step 4 2). Value = Stack [TOP] 2) TOP = TOP+1 3) TOP=TOP-1 Stack [TOP]=Value 2) 4) Stop A) Stop. peek: Als' 1). If TOP = NOLL Error "Stale is empty" so to step 2. 2) · Return Stack [TOP] 3) Stop. Program Code: #include Lstdio.h> # include < stillib.h> #include < conjunt> # include a control // Altering this value changes size of stack created. ent st[MAX], top=1; void push (int sl), int val); int pop (int st(3); int peek (cirt st[]); void display (int still) int main (int ang c, char & ang v []) int ral, option;

do printf ("In *** * Main Menu *** * "); 2 mintf ("In 1. PUSH In2. POP In3. PEEK IN4. DISPLAY print f ("In Enter the option: "); scenf ("y.d", & ophon"); Switch (option) case ". printf["In Enter the number to be pushed on store" scanf ("Id", & val); push (sty val) break; cane 2: val= pop(st); if (val !=-D pointf ("In the value deleted from stace is Y.d" break; care 2: val=peek (st); printfl" In The value stored at TOP is ' Yd', val); g case A: display (st); break; pubile (option := 5); 3 returo; roid push (int s[], int val) 2 if Ctop == MAX-D. printf ("In Stack Overflow");

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else toptt) st [top] = Val; int pop(int stip) int val; if (top==-1) printf ("In Stack Under flow"); 3 return -1; ehe 1. val = st [top]; y vetuin val; zvoid display (ent sc3) int i; obe for (;= top; i>=0; i=-) 31 mintf ["In 1. d', St[i]); 3 int peek (int st[]) 3 1f(top==-1) mintfl" In Stack is empty"); rehun -1; éhe retur (st[top]); 3

(8)

Linked List Implementation of Shalk;-

Type declaration;-# ifndet _Stack_h smut node; typedef shuch Noole & Pts To Node; typedef PtrToNode Stack; int Is Empty (Stace S); Stack Createstack (void); void Dispose Stack (Stack S?; roid Make Empty (Stack S?; Void Posh (Element Type X, Stack S); Element Type Top (Stack S); Void Pop (Stack S); #endif 1t _stack_ht1 1 & Place in implementation file +/ to stack implementation is a linked list with a header of struct Node Element Type Element; PhrTo Node Next; 3; EsEmpty Routine: int Is Empty (Stack S) return S-7 Next == NULL; Create on Emply stack: Stack CreateStack Croid) Stack SI S= malloc (size of (smut Node)); If (S== NULL) FatalEmir (" Out of Space !!! ");

Maxe Empty (3); z return si Void Make Empty (Stack S) Emor (" In Must me Greate Stack First"); if (S==NULL) else while (! Is Empty (0)) Pop(S); Push : roid Park (Element Type X, Staces) PtrTo Node Trup(ell) Top Cell=malloc (size of (struct Node)); if (Implell == NULL) Fartal Error (" Out of Space !!! "); else Truple 11 -> Element = X; ToupCell -> Next = S-> Next; 3 StNext = Tmplell; ElementType Top(Stace 5) Jop : if (! Is Empty (3)) return S-> Next -> Element; Error ("Emply stace"); y seturn o;

· Pop: void Pop (Stack S) PhrToNode FirstCell; if (Is Empty (S)) Error (" Empty Stack ")) ehe E FirstCell = S -> Next; S->Next = S-> Next ->Next 3 3 } face (First (ell); Representation: 210 1 72 74 7 Fig. Linked List representation of Stack TOP Here if TOP == NULL, then the stack is empty. The head pointer of the linked dist is used as TOP. All the insertions (punk and deletions (pop) are done at the node pointed by TOP. Opuations on Stack: D. Punk: >3 Fig. Before push of 15 79) TOP -+ - #3 JI Fig. After push of 15 29/7 15 -> 10 TOP Ale: 1). Allocate memory for the new node and name it as NEW_NODE 2). NEW_NODE -7. Data = Value. 3). If TOP = Null NEW_NODE -7 Next= Null TOP = NEW_NODE

Elese

NEW_NODE -> Next = TOP TOP = NEW_NODE

A). Stop.

2). Pop:

10 3 TOP Before deletion of TOP 3 4 10 After deletion TOP of TOP Als: If TOP = Null. print "Underflow ptr=TOP 2), 3). TOP = TOP - Next A) free (ptr) 5) Stop Program Code: findude Kstdio, h> #include L conio. h> #include <stdlib.h> # include < malloe. h> Smut stack 2 int data; smut stack & next; 3; Storet Stack + top = NULL; smut stack & push (shout stack +, int); struct stack to display (struct stack +); struct stack + pop(struct stack +); int peek (smich stalk +);

Int main (int aige, char taigres) 3 int val, option; do 5 promif ("In tot * + Mais Menu tot to dod"); print f ("In 1. PUSH IN 2. POP IN 3. PEFK IN4. DURAY 1n5. Ex17 "); print of (" In Enter your choices"); scenf ("Y. d" & option); switch (option) print f ("In Enter the number to be pashed:"); Lone 1: Scenf ("1.d", & val); top= push (top, val); break; care 2: top=pop(top); break; Case 3: val = peek (top); if (val != -1) printf["In The value at TOP is ".d", val); the pointf[" In Emply Stack "); preak', Cone A: top = display (top); break; Juhile (option! =5); return o;

Struct stack * push (struct stack * top, int val). 3 shut stack * ptr; ptr= (struct stack to) malloc (size of (struct stack)); phr->data =val; if (top== NULL) ptr-7 Next = NULL; top=ph; etre ptr->next = top; , top=ph; 3 return top; Smit stack & display (street stack & top) shut stack & ptr; ptr= top; if (top==NULL) print f C" Empty Stack"?; else while (ptr != NULL) I mint f ("In Y.d", ptr->idata); j phr=phr->next; retun top; to shult stack & pop(smilt stack \$ top) 3 smut stack + ptr; ph=top;

if (top == NULL) printf("In Stack Underflow"); else · top= top= next; printf ("In The value being deleted is Y.d", phr. >data); 2 Free (prr); z return top; int peek (struct stack top) if (top == NULL) ebe return top - data; Applications of Stack;-1). Expression Convolution (b) Infix to prefix (a) Infix to postfix (d) Prefix to infix (c) Pestfix to infix 2. Expression Evaluation Parsing 3), Similuation of recursion A). Function call 5). Balancing Symbols 6) Revousing a chit 7) Recurion A) Tower of Hanoi. 9)

Evaluation of Arithmetic Expressions: Expression Representation:

3 methods 1). Infix x+y operator bla operands 2). prefix +xy operator before operands 3). postfix scyt operator after operands Postfix: The expression (A+B) + C can be written as: [AB+] & C AB+ C & is the postfix notation. The postfix expression does not even follow the rules of the operator first in the expression is operated first the operands. of operator precedence. on the operands. I For example, given a postfix expression AB+C+. While evaluation, addition will be performed prior to multiplication. & This, operators are applied to the operands that are immediately left to them. The prefix expression is evaluated from left to right. Prefix: & Here, the operator is placed before the operands.

& Here, the operator in places acquired is fix is infix & For example, if Atts is an expression (isfix) is infix notation, then the corresponding expression is prefix notation is tAB. While evaluating a prefix expression, the notation are applied that are present immediately on the operators are applied that are present immediately on the right of the operato.

(16)

Conversion of an Infix into a Postfix Expression: Let I be an algebraic expression written is infix notation. I may contain parentheses, operando and operators. The precedence of these operators ben be given as: Higher priority \$, 1, %. Lowel priority t, -Ex: Infix to prefix 1) (A+B) &C 3) (A+B) / (C+D) - (D+E) [HAB] / (CD] - [ADE] (JAB)&C [1 +AB+CD] - [4 DE] \$ +AB - ITABTOD *DE 2) (A-B) * (C+D) [- AB] * (+CD] X-AB+CD Infix to postfix 2), (A+B) / (C+D) - (D*E) 1) (A-B)*(C+D) [AR+] / (CD+] - (DE+] [AB-]*(CD+) [AB+ CD+]-[DE+] AB-CD+ & AB+CD+/DE+-* The order of evaluation of these operators can be charged by making we of parenthesis. * For example, (A+B) *C, will evaluate A+B front and the the result will be multiplied with C, The alg accepts an infix expression that may contain operators, operands, and parenthesis. of the alg uses a stack to temportanily hold operators. * The past fix expression is obtained from left to -right ming the operands from the infix expression, and the operators which are removed from the stack.

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the stack and to add a corresponding right parenthesis on the stack and to add a corresponding right parenthesis at the end of infix expression. * The alg is repeated with the stack is empty. 1). Add (Push) ." (" on to the stace 2). Repeat until each character in infix notation is scanned if a "(" is encountered, push it on the stack if an operand (whether a digit or a character) is encountered, add it to the past fix expression. if a ")" is encountered then (a). Repeatedly pop from stack and add it to the pastfix expressionentil a "C" in encountered. (b). Discoold the "(". That is, remove the "(" from stack and do not add it to the post fix expression. if an operator O is encountered, then (a). Repeatedly pop from stack and each operator (popped from the stace) to the postfix expression which has the same precedence or a higher precedence than 0. (b) Push the operator O too the stace, 3). Repeatedly pop from stace and add it to postfix expression until the stace is empty.

A) EXIT

x Character	- (B/C + (D7. E F Stack	Post fix Expression
canned	C	
	(A
A	C -	A
15. 1. 1.	1-1	A
C	(C	AB
B	L	AB
/	C - C/	ABC
C	C- (/	ABC/
+	€- (+	ABC/
((-(+(ABC/D
D	(- (+ (ABCID
1.	C- (+ ()	
12	6- (+(%	
*	(-(+C %	1000
F	(- (+())	ABC DEF
1	(-(+	ABCIDEF of 7.
)	C- C+/	ABC/DEF + %.
1	(- (+/	ABC/DEF&% G
G		ABC / DEF& /. GI+
)	(-	ABC (DEF& /- Q1/+
¢	(-* (-*	ABC/DEF* / GIL+H

Evaluation of Post fix Expression; -

Ale 1). Add a ")" at the end of the pastfix expression 1). Scan every character of the post fix expression and repeat step 3 and 4 with ")" is encountered. 3). If an operand is encountered, push it on the stace. If an operator O is encountered, then (a). Pop the top two elements from the stack as A and B as A and B. (b). Evaluate BOA, where A is the topmost element and B is the element below A. (c). Push the result of evaluation on the stack. 4). Set Result equal to the topmest element of the stack. 5). Exil Example: Postfix Expression 9- (3+4)+8)/4. - = 934 # 8+4/postfix Stack Character Scanned 9 9 9,3 3 -9,3,4 4 9,12 * 9,12,8 8 9,20 + 9,20,6 4 9,5 1 4 Here, in fix expression is converted into post fix from then

this alg is applied.

Pam. Code to Conversion: Infix to Postfix

include < stdio. h> #include 2 conio. h #include < ctype. b? # Include a string. h # define MAX 100 char st[MAX]; int top=-1; void push (chan st[], chan); char pop (char st[]); void Infix to Postfix (char source 13, char taget [3); int getPriority (choor); int main () char infix [100], pastfix [100]; printfl" In Enter the infix expression: ">; gets (in fix); shopy (postfix , " "); InfixtoPostfix (cofix, postfix); printf("In The corresponding postfix expression is"); pith (postfix); z retuno; Void Infix to Pestfix (char source [], char target []) int 1=0, j=0; char temp; stropy (target, " "); while (source [i] != '10') if (source Ei] = = '(') push (st, source [i]); z itt;

22 else if (source[i] == ')') while ((top ! = -1) && (s+[top]! = '(')) 3 target [j] = pop(st); if(top = = -i)2 printf["In Incorrect expression.")) z exit(1); temp = pop(st); Premore left parenthesis & i++; else if (indigit (source [;]) || isolpha (source [i]) target[j] = source[i]; else if (source [i] == '+' || source [i] == '-' || source [i]== '+' 11 source[i] == '1' || source[i] == '1') 5 while ((top ! = -1) && (st [top]! = '(') && (get "in only (st (top)) > get Priority (source [i))) target[j] = pop(st); 3 Jt+; punh (st, source [;]); i++;

ehe printf (" In Incorrect element is expression."); zexit(1); while ((top! = -1) & (st[top]! = '(')) target (i) = popest); 20++; , touget[]] = '10'; Int get Priority (char op) if (op=='/' ll op== '*'llop=='/') else if (op == '+' || op == '-') 3 return o; void push (chas st(3, chas val) if (top == MAX-1) print f["In Stack Overflow !!! "); else topt; 3 - 3 st[top] = val; char pop(char stc]) ? chou val = ' printf("In Stace Underflow !!! "); if Crop = = -1/2else val = st[top]; top --; Ireturn val;

Jam Cade for Evaluation of Partfix Expression:

Atinelinde estelie.hz Hindude & conishs #Include ectype.hy Adaptine Mar 100 float stimdx7; push lafloat stl7, flor val); Part top = -12 vald per (floar st 3); flead evaluate Portfix (char exp(3); fleat int main() flot vals chai explices; printf(" In Enter the postfix expression:"); printf[" In Value of Postfix expression = 1. 20 f", val); gets (exp); return o; float evaluate Pastfix (char exp[]) int i = 0; fleat op!, opa, value; while (exp[i] = 10) else push (st, float) (exp[i] - 'o')); if (isdigit (exp(iD) opa = pop(s+2; opia poplat);

Igm Code for Evaluation of Past fix Expression continuation. switch (exp[i]) Case value = op1+op2; break; Couse Value = op1- op2; break) Carl Value = opi/opa; break; Care value = op1 + op2; break; value = (int) op! 1. (int) op?; Conc g. push(st, rahue); return (pop(st)) puch (float st[], float val) void 2 if (top==Max-D printf("In Stack Overflow !!! "); she top ++; st[top] = val; 3 floot Pop(float st(3) float val = -1; if (top==-1) printf (" In State Underflow !!! "); else val= st[fop]; top--; 9 return val; 3

Conversion of Infix to Prefix Expression:

Alg 1: 1). Reverse the infix string. Note that while reversing the string you must interchange left and right 2). Obtain the postfix expression of the infix expression parentheses. 3). Revenue the postfix expression to get the prefix expression. Example: Infix expression=7 (A-B/C) + (A/k-L) Step!: (L- K/A) of (C/B-A) - Reverse the string Step 2: Postfix expression of (L-K/A) & (C/B-A) is got as [LKA1-] + [CB/A-] LKA 1- CBIA- & *-A/BC-/AKL =) Reven the postfix shing. Step 2: 1). Scan each character in the infix expression. For this, repeat Alga Steps-8 until the end of infix expression. 2). Push the operator into the operand stack, and ignore all the left parenthesis until a right parenthesis is encountered 1). Pop operand & from operand stack 4). Pop operand 1 " 5). Pop operator from operator stack. Concatenate operator and operand. Concatenate result with operand?. Push result into the operand stack. 6). 7), 8), 9) END.

Pam Code to convert infix to prefix expression:

#include < stdio. h7 # include < conio.h7 # include a string. h7 # include < ctype . h> # define MAX 100 char St[MAX]; int top=-1; Void revenue (char strE J); push (char sf[], char); Void shat pop (char s+(3); void Infix to Prefix (char source [], char target []); int get Priority (chan); chan infix [100], postfix [100], temp[100]; int main () the infix expression: "); 3 CYADO; printf ("In Enter gets (infix ?, revouse (infix); Stropy (postfix, ""); Infix to Prefix (temp, postfix) postfix expression is: "); printf(" The corresponding put (post fix); stropy (temp, ""); reverse (postfix); printf["In The prefix expression is :"); puts (temp); 3 getun o;

+aviget
$$[j] = pop(st);$$

 $j+t;$
 $j \in (top = = -1)$
 $print f(' |n Incorrect expression');$
 $exit(D;$
 $femp = pop(st); //remove left pacentheris$
 $j+t;$
 $j = he if(iadigit(source[i]) || isolpha(source[i]))$
 $fauget[j] = source[i];$
 $j+t;$
 $j = he if(source[i] = = '+' || source[i] = = '-' ||$
 $source[i] = = '+' || source[i] = = '-' ||$
 $source[i] = = '+' || source[i] = = '/' ||$
 $source[i] = = '+' || source[i] = = '/' ||$
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 $source[i] = = '+' || source[i] = = '/' || source[i] = : '/' ||$

30 else printf ("In Incorrect element is expression."); } exit(i); ²while ((top!=-1) && (st [top]!=`(')) target [j] = pop(s+); } target[j]= `\o'; 3 jut get Priority (char op) if (op=='/' || op== 'd' || op== '1-) else if (op== '+' || op== '-') retuno; void punk (char st[], char val) { if Ctop == MAX-1) printf ("In stace Overflow"); else toptt; sfEtop]=val; 2

chai pop (chan str]) char vals'; printf ("In Stack Underflow"); if (top==-) else val =st[top]; 3 top - - 1 3 return val; Evaluation of Postfix Expression;-1) Accept the proefix expression Als: 2) Repeat until all the characters is the prefix expression have been scanned. (a) Scan the prefix expression from right, one (b) If the scanned character is an operand, push it on the operand stack. (c) If the scanned character is an operator, this (i) Pop two values from operand stack (ii) Apply the operator on the popped operands (ii) Push the result on the operand stack. Prefix expression +-927 + 8/4 2. Example:

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character.	Operand Stark
scanned	Stalk
12	12
4	12,4
1	3
8	3,8
*	24
7	24,7
2	24,7,2
-	24,5
t	29

Pen code to evaluate prefix expression: #include < Stdio. h> #include < conio.h> #include <shing. 17 #include < ctype. h> int st E107; int top=-1; int popes; void push (int); int main) E chas prefix [10]; int len, val, i, opri, opra, res; chracici printf [" In Enter the prefix expression:"); gets (prefix); len = styles (prefix); for (i = len - 1; i>=0; i-=)

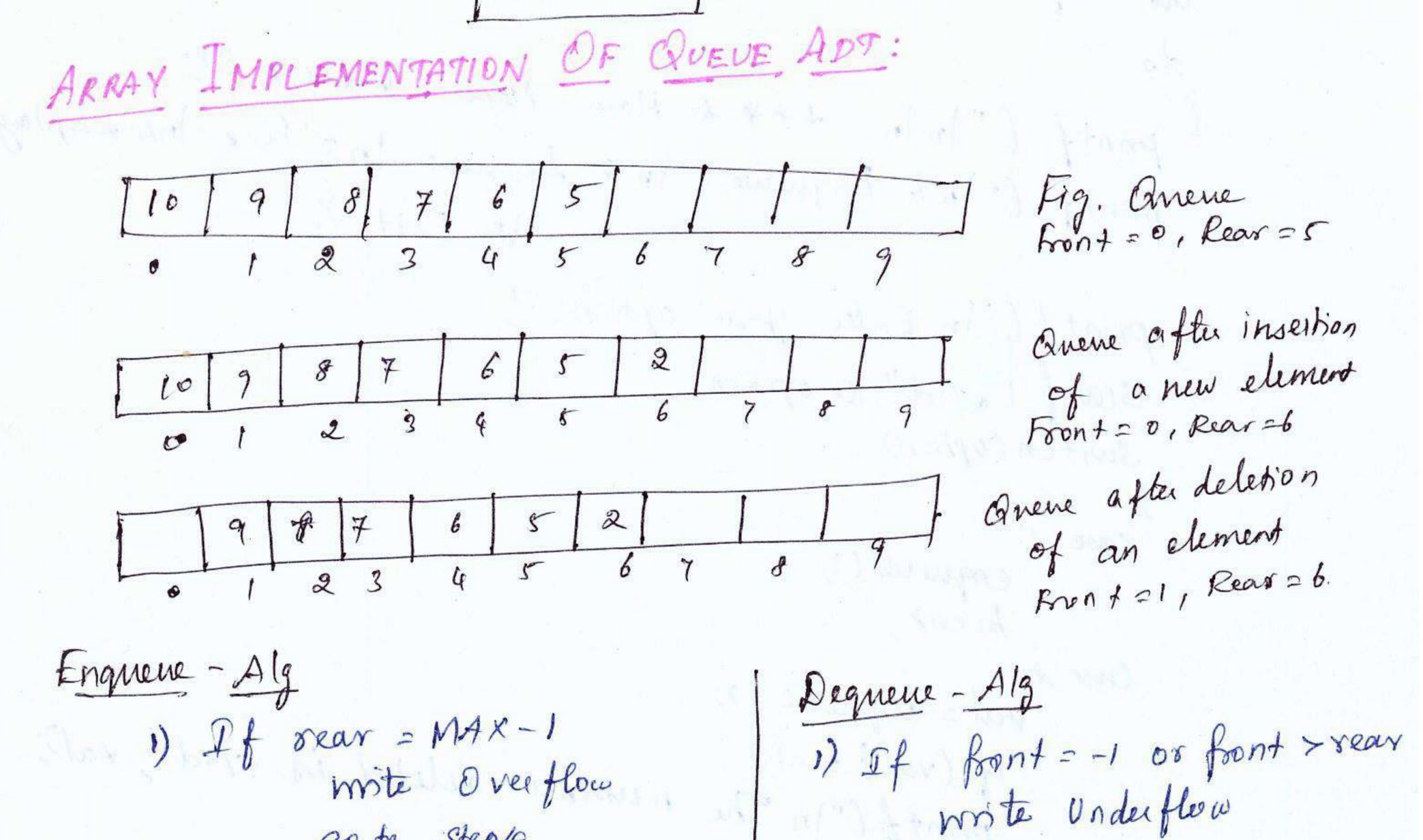
Switch (get type (prefix[i])) 3 Case o: val= prefix[i] - '0'; push (val); break; Cone 1: opri=popli; opra=pop(); switch(prefix[i]) 5 case '+': res = opr1 + opr2; break; case '- ': res = opr1 - opr2; break; Case ' * : res = opr1 * opr2; break; Case '1': res= opr1/opr2; break; } push(res); printf ("In Result = Y.d", st[0]); zeturno; int get_type(charc) if(c=='+' // c=='-' // c=='+' void proh(int val) L st[++top]= val; Il C == '/') return 1; else return 0; z 3 int pop() E return (st [top-]); Z



Queue is a linear data structure with First In, First Quit manner. (FIFO). The elements that inserted first is the first one to be taken out. The elements in a gnene are added at one end called the <u>REAR</u> and removed from the other end called the <u>FRONT</u>.

The basic operations on a gneve are Engueve, which inserts an element at the end of the list (sears), and Degneve, which deletes the element at the start of the list (front)

Dequeue (Q) Avene Q Enqueue (Q)



go to Step4 a) If front=-1 and rear-1 set front=rears=0 else set rear= reart, END If 3) Set Queue [rear]=num A) EXIT

else set val = Queue[foont] set front = front ti End If &) EXIT

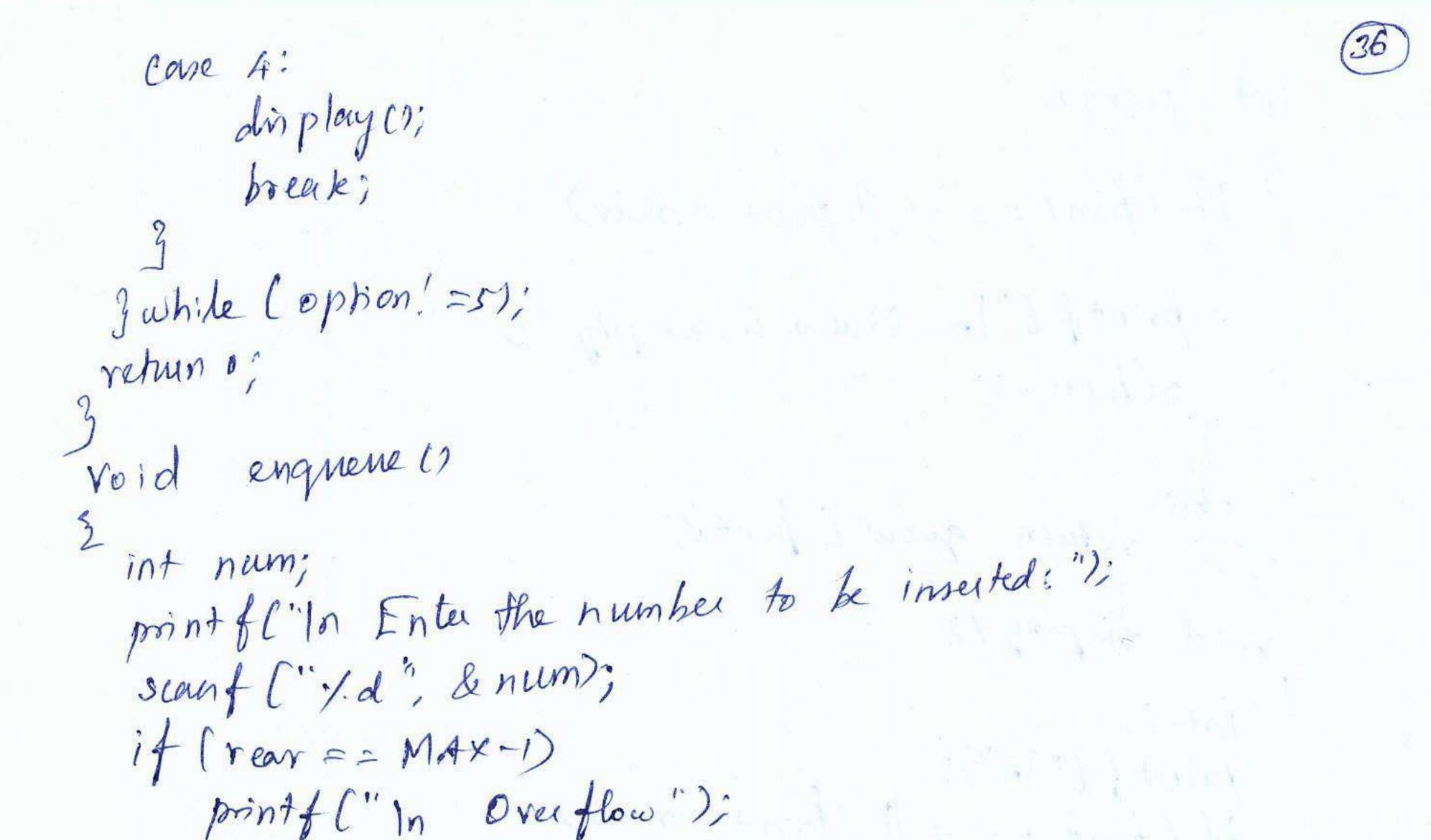
13m code:

l'Array Implementation Chure ADT Df #include < stdio. h> #include < conio. h> # define MAX 10 int queue [MAX]; int front = -1, rear = -1; the second second void inqueue (void); int degnene (void); int peek (void); void display (void); int main () int option, val;

do 2 printf ("In In + * * Mais Menu # * * * "); print f ("In 1. Enquence In 2. Dequere In 3. Peek In4: Display WE. Exit");

print f l'In Enter your option: "?; scanf ["1.d", & opsion?; Switch Coppion) case 1: enqueue (); boeak;

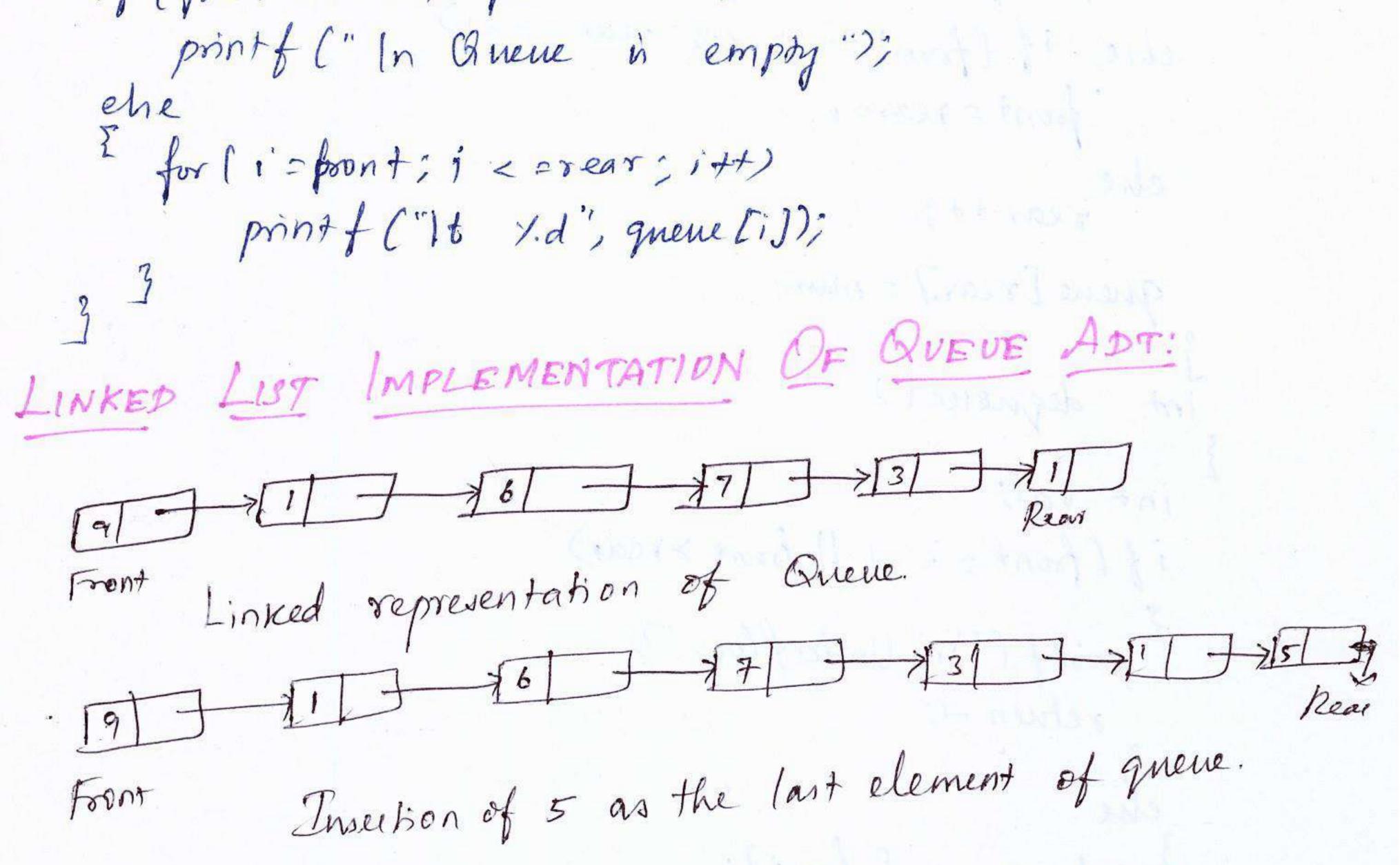
Come 2: val = degneire (); number deleted is : ".d", val); if (val! = -1) printf ("In The break; Cone 3: val=peek(); printf ("InThe first value is gneue is : Y.d", val); break;



else if (front ==-) c& rear ==-) foont = rear=0; else rear ++; guene [rear] = num; int dequene() int val; if (foont ==-1 || foont > rear) if (foont ==-1 || foont > rear) int f ("In Underflow"); return -1; }

val= queue [foont]; foont ++; if (foont > read) foont = read = -1; return val;

int peek () if (front = = -1 || front > rear) printfl"In Queue is empty"?" zehun -1; z che return grune [font]; void display () int ig printf("In"); if (front == -1 1) foort > rear)



Deletion of 9 as the first element of queue. Front

The strage requirement of linked representation of a queue with a element in O(n) and typical time for openation in O(i). In a linked queue, every element has two parts-one that stores the data and another that stores the address of the next element. A the START pointer of the linked list is used as FRONT. A Here, we will also use the another pointer called REAR, which will store the address of the last element in the mene. A All insertions will be done at the rearis end and

all the deleksons will be done at the front end. #If FRONT=REAR=NULL, then it indicates that the queue is empty. Inselt (Enqueue) Alg: 0) Allocate memory for the new node and name it co pTR. 2) Set PTR = NULL 3) If FRONT = NULL Set FRONT = REAR = PTR Set FRONT = NEXT = NULL

else REAR-INEXT = PTR Set REAR=PTR Set (End if) A ENA REAR =>NEXT=NULL

Deletion (Degnewe) Alg: ?) If FRONT = NULL Write "Underflow"

END IFJ 2) Set PTR=FRONT 3) Set FRONT = FRONT -7 NEXT 4) Free PTR 5) END and such they are such Pan Code: mining and Le la maine all ante ditta addition finclude < stdio.h7 finclude < comio. h> finclude < malloch7 Smet node int data; Spuct node thert; Smit queue Spuct node & front; Spuct node & rear; 3; smit quere +9; void create - queue (struct queue x); struct queue dinsent (sput queue 2, in2); struct queue *delete element (struct queue *);

struct quine # display (shuct quene #). int peer (struct quene t); int main () int val, oppon; create_queue(q); chancy,

printf("In *** Main Menn **** "> mn+f["In 1. Inset In 2. Delete In 3. Peer In4. Diplay printf("In Enter your option:"); scanf("'/d", &option); switch (option) to insect in gnene:"? Couse 1: printf ["In Enter the no. Scanf (". Y.d", &val); q= insert (q, val); preak; 2 = delete-element (9); Cone 2: break; Cane 3: val = peek(q);if(val = = -i)pointf("In The value at front of queue is '.d", val); break; Case A: q=display(q); break; getcher; returno;

Void create-queue (struct queue +q)
2
2-7reax = NULL;
2 -> foont = NULL;
struct queue * insent (struct queue *q, int val)
2
Struct node * ptr;
·ptr = (struct node *) malloc (size of (struct nede));
ptr -> data = val;
if (q-7 foont = = NULL)
2
q > foont = ptr;

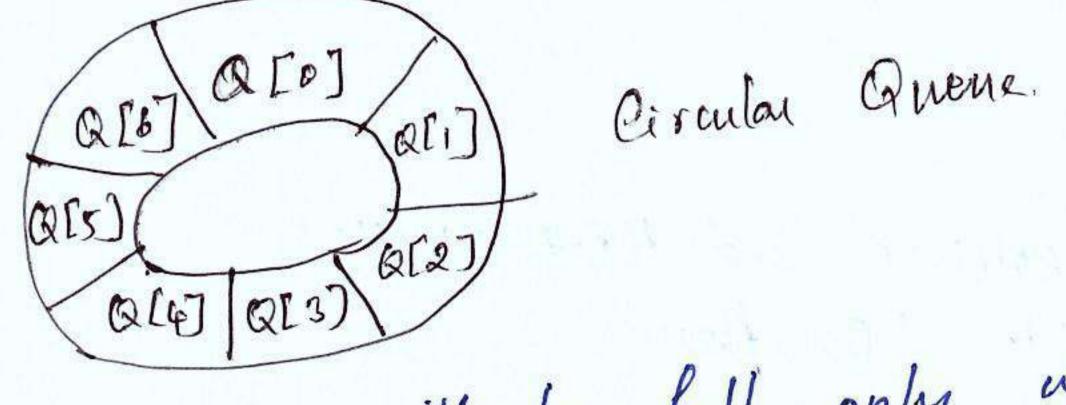
2-7 rear = ptr; 2-7 front-inext = 9-irear -inext = WULL; ene 2-) rear - Thext = ptr. 2-7 rear = ptr; 2-7 rear -7 next = NULL; return 9; stonet queue & display (stanct queue *q) somet node appr; ptr=27 foont; if (ptr = = NOLL) printf("In Queue in empty")

else printf("In"); while (p&:!= 2-7 rear) 2 printf ("',d It", ptr-7 data); 2 pts= ptr-7 next; printf (" 'Alt," phr-7 data); retunz; Strict queue #delete_element (strict queue *q) struct node & ptr; pt= 27 font; if (2-) front = = NULL) printfl" In Onderflow"); else 9-7 font = 9-7 front -7 next; printf("In The value being deleted is 7.d", ptr-7 do por-7 data); free (por); return qi int peer (struct greve #2) if (2-) front = = NULL) 2 mintf ("In Queue is empty"); setun -1; zebe return 2) foont -> data;

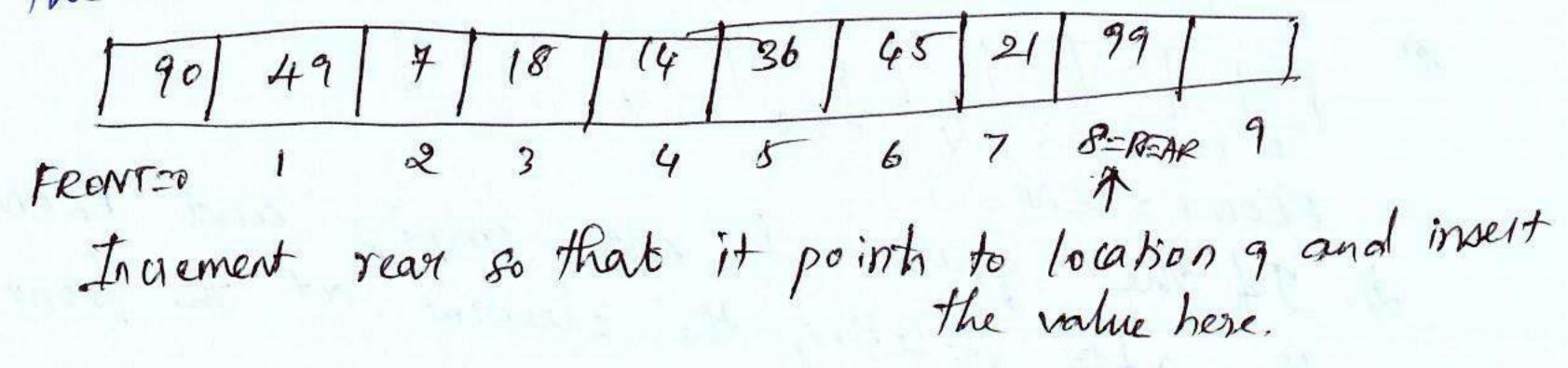
(43) Types of Queues; -J. Circular Queue 2). Deque (Double Ended Queue) 3). Priority Queue 4), Multiple Anene STATE A STATE Freder - proved - short Circular Queue

In linear gnewes, insertion can be done only at one end called the <u>REAR</u> and deletion are always done at the other end called the FRONT. 549718143614512199172 Hue FRONT = 0, REAR = 9. Now, if you want to inselt another value, it will not be possible because the gnene is completely fall. Thue is no empty space where the value can be inserted. Comider a scaravio in which two successive quere will then be given as de le kons ave made. The - 7/18/14/36/45/21/29/72 0 / 23 4 5 6 7 # 9 Chuene after two successive deletsom. strown below: Here FRONT = 2, REAR = 9. Suppose we want to insert a new element in the quene. Even though there is space available, the overflow condition still exists because the condition REAR = MAX-1 Still hold tone. This is a major drawback of a linear quene.

To resolve this pbm, we have @ solution: 1) Shift the elements to the left so that the vacant spaces can be occupied and utilized efficiently. But this can be very time-consuming, especially with when the gneve is large. 2) Using Circular Queue, the first index comes after the Cart index.



The circular queue will be fall only when FRONT=0& & The circular gnere is implemented in the same manner REAR = MAX-1. to the only difference will be in the code that performs insertion and deletion operation. Insertion of for insertion, we have to check the @ condition: as a linear queue is implemented. i) If FRONT=0 and REAR= MAX-1, then the circular quere à full. 90 49 7 18 64 36 45 21 99 72 Full Quene. FRONT=0 1 2 3 4 5 6 7 8 REAR=9. 2) If REAR! = MAX-1, then REAR will be incremented and the value will be inserted as



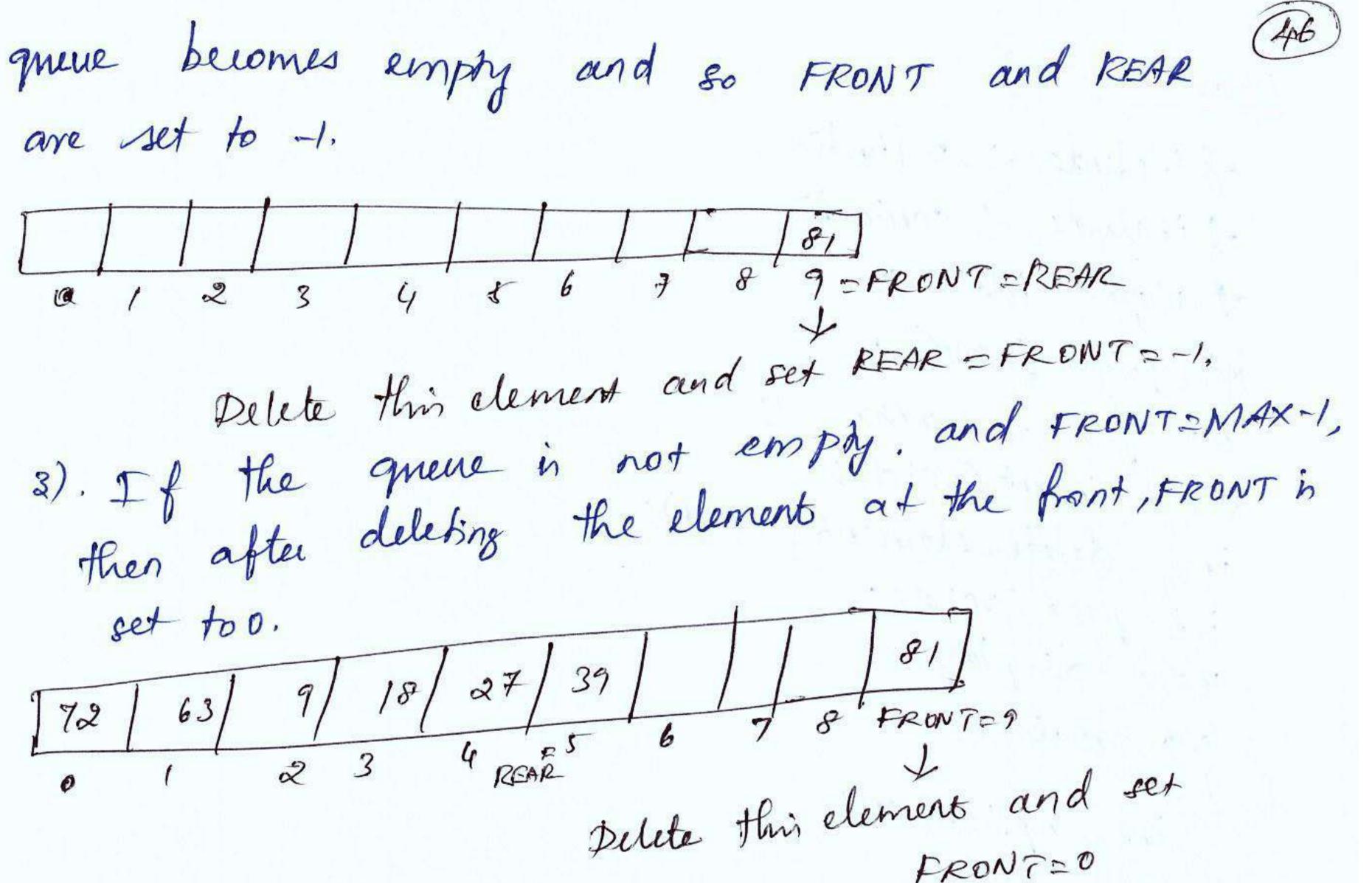
3) If FRONT!-0 and REAR = MAX-1, then it means that the queue is not full. So set REAR = 0 and insert the new element there as in Fig. below.

7/18/14/36/45/21/80 P1 of 'FRONTE2345678 REARON Set REAR =0 and inself the value here.

Algorithm:

1) If FRONT=0 and REAR=MAX-1 Write "Overflow" goto step 4 END IL 2) If FRONT = -1 and REAR = -1 Set FRONT = REAR = 0. Else If REAR = MAX-1 and FRONT! =0 A LA Set REAR=0 Else Set REAR = REARTI ALL AND ALL AND ALLAND the sumption and such that is a ENDIF 3) Set QUEVE[REAR] = VAL 1.612.84 A) FXIT To delete an element from circular gneve, we check B To delete an element from circular gneve, we condition: Delebion: 1) Look at Fig below If FRONT --- , then there are no elements is the gnewe. So, an cenderflow condition will be repeated.

3 a). If the gnine is not empty and FRONT= REAR, then after deleting the element at the front the



FRONTED

Alsonithm:) If FRONT =- 1 Write "Underflow" Soto steps END If 2) SUT VAL = QUEUE [FRONT] 3) If FRONT = REAR Sed FRONT = REAR-1 Else If FRONT=MAX-1 Set FRONT=09 Set FRONT = FRONTH Else END of





Pgm Lode: # include < stdio.h> # include < conio.h> # define MAX 10 int queue [MAX]; wit font = 1, rear = -1; void inselt (void); int delete element (void); int delete element (void); int peek (void); void din play (); void din play (); int option, val;

chaser (); printf ("In d ## & Mais Menu & # & # "); printf ("In 1. Inselt Ing. Delete Inz. Peek In4. Dis play In 5. Exit"); print f (" In Enter your option!"); scarf (" Y. d", & option); switch (option) case l: inverto; break(); val = delete element (); care 2: pointfl" In The element deleted in ".d" if(val==-) break; Case 3! val= peck(); if(val==-1) if(val==-1) pointf("In the fixet value in v.d", val); break;

Case 4: dis play (); break; 1111 Jubile (option!=5); XX j return oj void insert() printf ("In Enter the number to be inserted in the queues"); Scanf (" '.d', & num); . >





 $i t \ l \ \beta ont == N$ foont = 0; ehe $foont \ t \ i$ y return val; int peck() ? if (front == -1 && rear == -1) E print f ("In Queue is empty"); return -1; the return gneue [front];

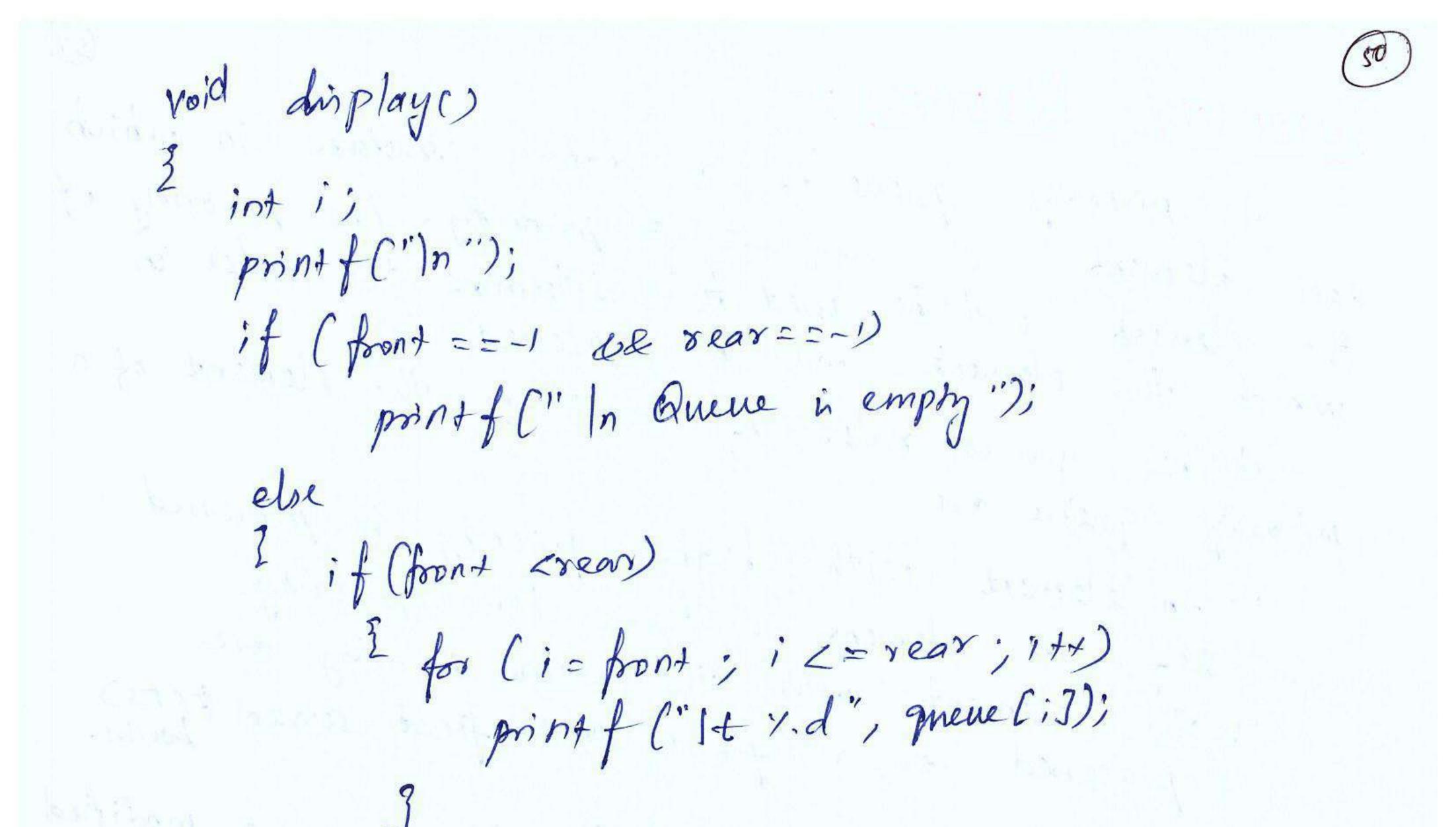
5.8

3 81 37 18

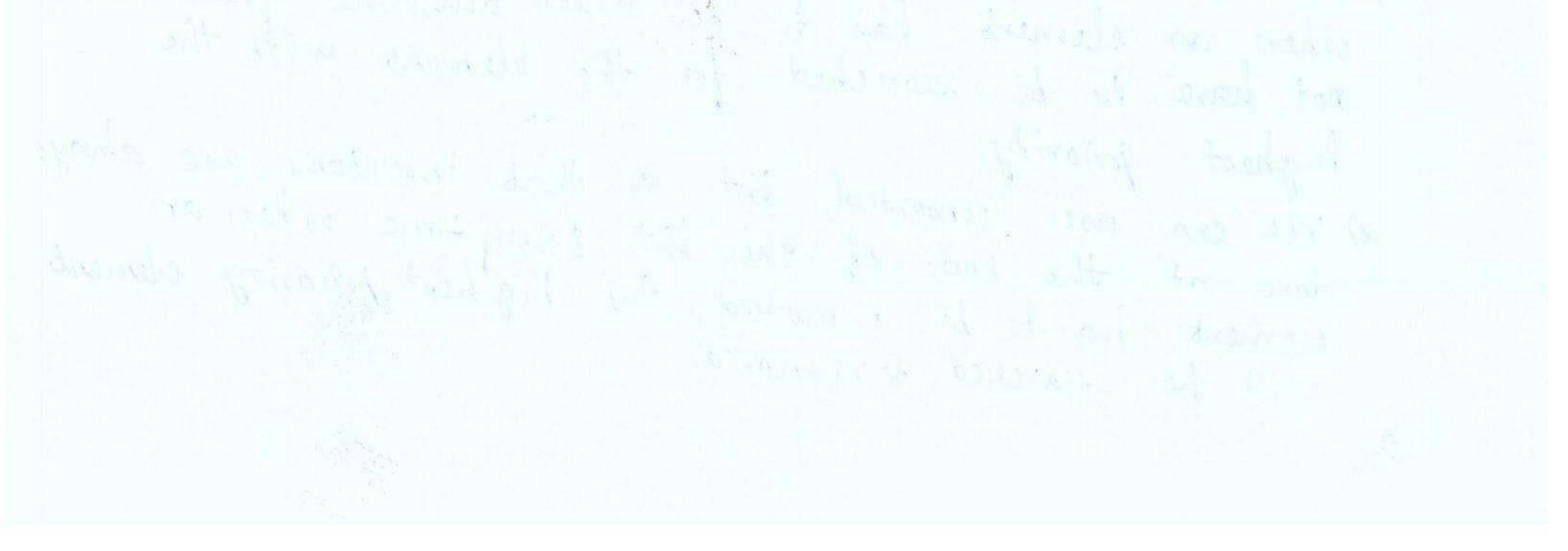
All and have

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elie 1 for (i= front ; i < MAX ; i+t) printf(" 1t y.d", queue [i]); for (i=0', i <= rear; i+t) printf(" \t y.d", queue [i]); g g g g



PRIORITY QUEUES:-A priority queue is a data structure in which each element is assigned a priority. The priority of the elements will be used to determine the order in which the elements will be proceed. thick the general rules of processing the elements of a priority queue are. 1. An element with higher priority in processed lefore as elements with a lower priority. 2. Two elements with the same priority are proceed on a first-come - first-served tects banis.

& A priority gnene can be thought of as a modified queue in which an element has to be removed from the queue, the one with the highest - priority is retrieved first.

to There are widely med in operating systems to execute the highest priority process filest. If the priority of the process many be set bound on the CPO time it requires to get executed completely. Implementation of a Priority Queue: (2) ways to implement a priority queue. (3) ways to implement a priority queue. (3) We can use a sorted bit to store the elements so that when an element has to be taken out, the queue will not have to be searched for the elements with the highest priority. (3) We can use unvorted birt so that insertions are always done at the end of the birt. Every time when an element has to be removed, the highest priority elements will be searched & removed.



A sorted list takes O(n) time to insert an element
in the list.
If takes only O(1) time to delete an element.
* An unsorted list will take O(1) time to insert an element
and O(n) time to delete on element from the list.
a blend of these two approaches in adopted that takes.
Dinked Representation of Priority Queue;
When a priority queue is implemented ming a linked his.

every node of the bit consists of 3 parts 1) the information I data part 2) the priority number of an element # If we use a sorted linked hist, then the element with The higher priority will precede the element with lower postonity. ATT BR + C3 - D3 - Flax dava mioning no Prioning Queue. In the above priority queue, lower priority number means higher priority. For example, if there are two elements A and B, where A has a priority number 1 and B has a priority no &,

where A has a processed before B. then A will be processed before B. I when two elements have the same priority the elements are orranged and processed in FCFS principle. There time the Served.

(53) Insertion: For insertion of an element, we have to traverse the entire list until we find a node with lower priority than that of the new element is inserted before the node with the lower prioring. & Suppose, if two elements has the same priority, the new element is inserted after that element. TATIF AB 2 Fols Jols JEGG (i) ATT TRE CI3 FIFT D 5 F JE/6/-J After insultion of F with portonity A. (ii) ATT JB2 JAT F2 JOI JAN (in After insultion of F with prioring 2 (same as B) After insultion of F with prioring 2 (same as B) Delesson is very simple process, where the first node of the list will be deleted and the data of that node will be processed first. Array Implementation of a Priority Queue: When arrays are used to implement a priority queue, then a separate queue for each priority number is maintained. & Each of these gnenes will be implemented ming circular arrays or circular queues. rEvery individual queue will have its own FRONT Dana coursean in the and REAR pointers. Brite States



We use a two-dimensional among where each mene will be allocated the same amount of space.

FRONT	REAR	1	2	3 4 5	
3	3	1		A	F. D. ita M.
1	3	2 3	CI	>	Fig. Priority Queue Matrix
4	5	3		EF	Mabix
4	1	$4 \lfloor I$		GH	

FRONT[k] and REAR[x] contain the front and rear values of row k, where k is the priority number.

Insertion:

with priority k', add the element

I in it a new element with position	
To inself a new element with prompt ,	as
at the rear end of row k, while k a the	
the priority no of that element to inself on element R'	
The priority no of man to inself on elements in	
the priority no of that elements the priority no of that elements the priority 's', then the priority queue is the won below with priority 's', then the priority queue is the won below	
in the wordy grene	
with prioring 3 then i	
FRONT REAR 12345	
FRONT REAR	
FRONT 1 -A Fig. Priority queue	
1 3 2 BCD matrix after 4 1 3 R EF Insertion of a new 4 1 4 I GH element	
2 R EF inselsion of a new	
4 GH element	
4 1 1 412	
Delebion: De lind non-empty queue	-
The delate an element, first we find	ah.
to device in a doment of the first homen	19
Deletion: To delete an element, first we find non-empty gnewe To delete an element, first we find non-empty gnewe and then moleus the front element of the first non-em	

quere. quere in own assumption, the first non-empty quere to the one with priority no i', so A will be deleted and processed first Pgm code for Linked Implementation of Priority Queue: (5)

#include <stdio.hz #include < malloc. h> finclude < conjo.h7 struct node int data; int priority; smict node Anext; node & start = NULL; node finsert (struct node *); smit node & delete (smit node *); Stout display (struct node *); Smit void int main() int option; elascrer; print & C"In ### Mais Menu & # "); do print f["In I. Inselt Ind. Delete 103. Display In4. Exit"); printf ('In Enter your option:"); scant (" " scanfl" y.d", & option ?; switch (option) stant = inselt (stant); case 1: break; Conr 2!

stant = delete (stant); break; Case 3: display (stant); break; gwhile (option (=4);



shuct node & invert (shuct node & shart) E int val, pri; shut node & ptr, & p; pbr= (shuct node +) malloc [size of (shut node)); printf("In Enter the value and its priority:"); scanf(" y.d y.d", & val, & pri); ptr-7 data = val; ptr-7 priority = pri; if (start = NULL 11 pri < start -7 priority) 2 ptr-7 next = start;

* spriet node #1 if (start = = NULL) 1 printf("In Under flow"); return; 3

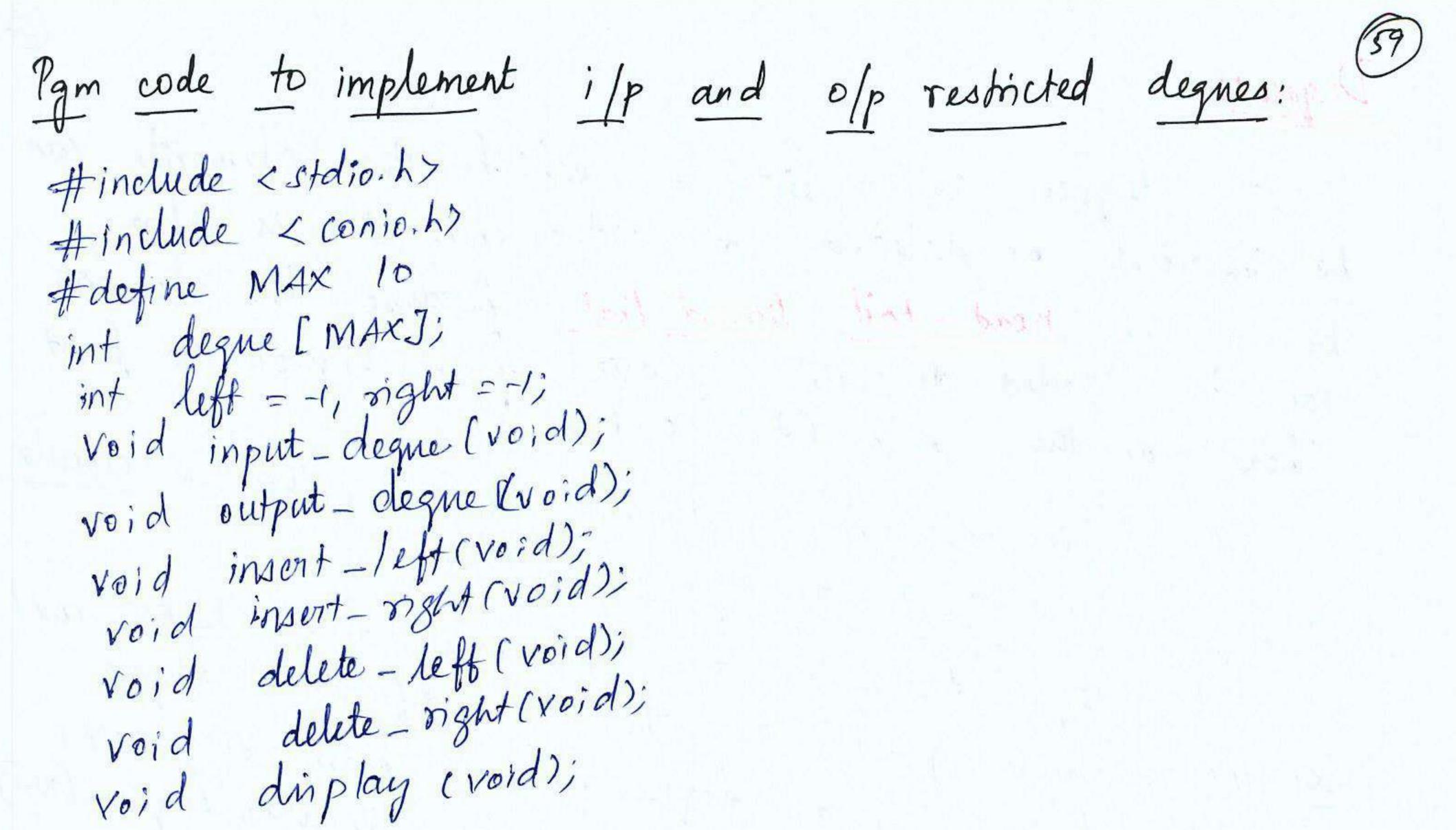






A deque is a list in which the elements can be inserted or deleted at either end. It is also known as <u>head-tail linked list</u> because the elements can be added to or removed from either the front (head) or the back (tail) end. A deque can be implemented using either a <u>circular</u> array or a <u>circular doubly linked list</u>. A deque, two pointers are mainstained, <u>LEFT</u> and of In deque, two pointers are mainstained, <u>LEFT</u> and RIGHT, which point to either end of the deque. RIGHT, which point to either end of the LEFT

a). Output restricted deque: - Deletions can be done only at one of the ends, while insertions can be done on both ends.



int main1? dated Mais Menu date"); 2 int option; print f("In 1. Input restricted deque"); print f("In 2. Output restricted deque); chacres; printf ("In Enter your option:"); scanf ("1.d", & option); switch (option) Case 1: input-deque (); break; case &! output-degne (); break; setur of



void input-degne ()

int option; do

printf("), INPUT RESTRICTED DEQUE"); printf("), 1. Insert at right"); printf("), a. Delete from left"); printf("), a. Delete from right"); printf("), a. Delete from right"); printf("), a. Display In 5. Quit"); printf("), Enter your option:"); Starf("'.d", & option; Switch (option)

1 case 1: insert - right(); break; case 2'delete -left(); break; case 3: delete - right(); break; case 3: delete - right(); break; case 4'display(); break; gwhile (option !=5); 3 void output - deque 1)

² int oppion; do ² printf ("In OUTPUT RESTRICTED DEQUE"); printf ("In I. Insert at right");

printf ("In 2 Insert at left"); printf ("In 3. Delete from left"); printf ("In 4. Diplay In 5. Quit"); printf ("In Enter your option:"); scarf ("It x. d", & option); Switch (option) 2 Case 1: insert- night(); break; Case 2: insert- left(); break;

Cone 32 delete - left (); break; case 4: display(); break; Zwhile Coption! = 5); void insert-night() int val; printf l'In Enter the value to be added: "); Scent ["1.d", & val); if [[left == 0 && night == MAX-D] [[left == night+1)] printf ("In Diverflow"); veturn;

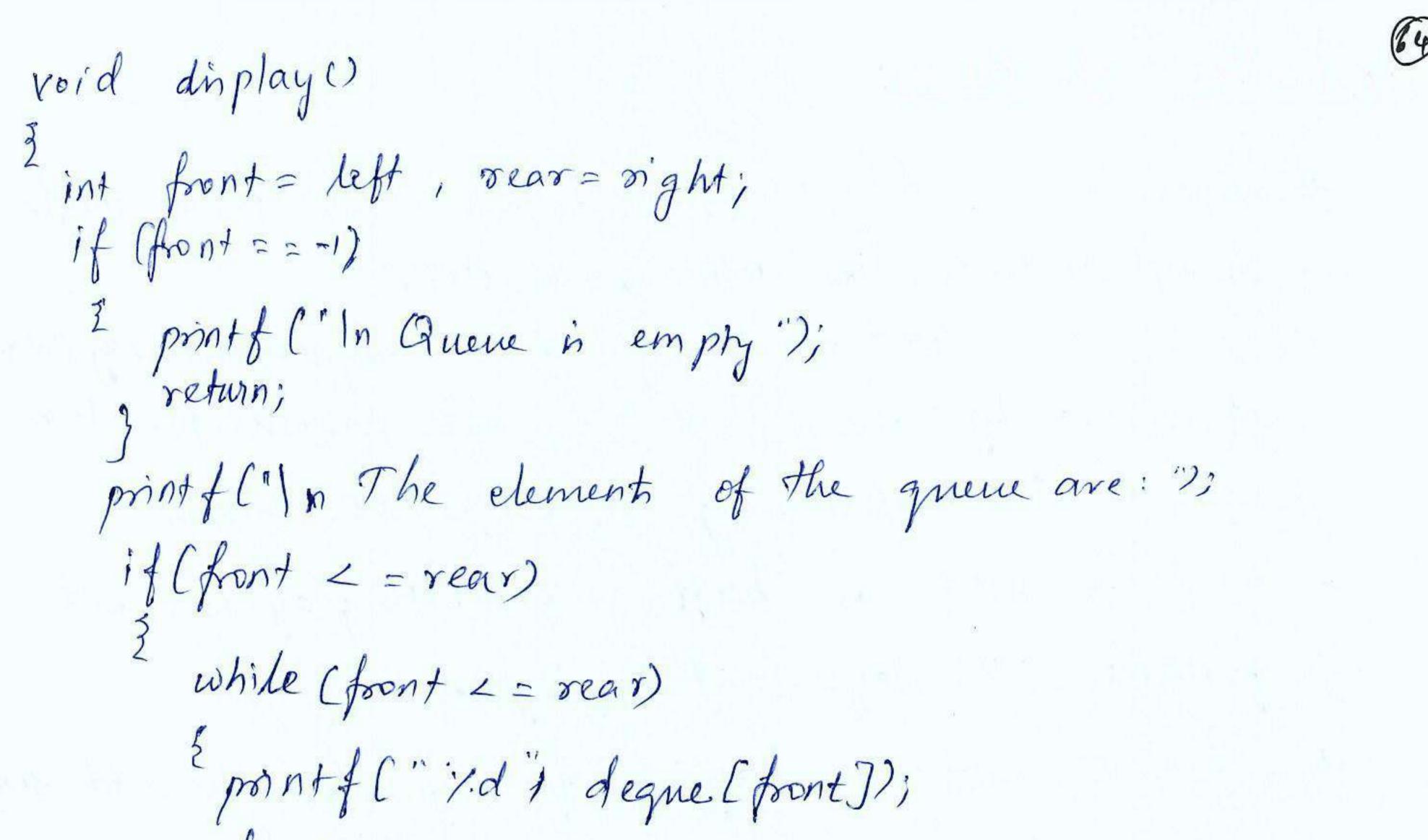
61

deque [left] = val;

V

63

ebe night = MAX-1; night = right -1; 0

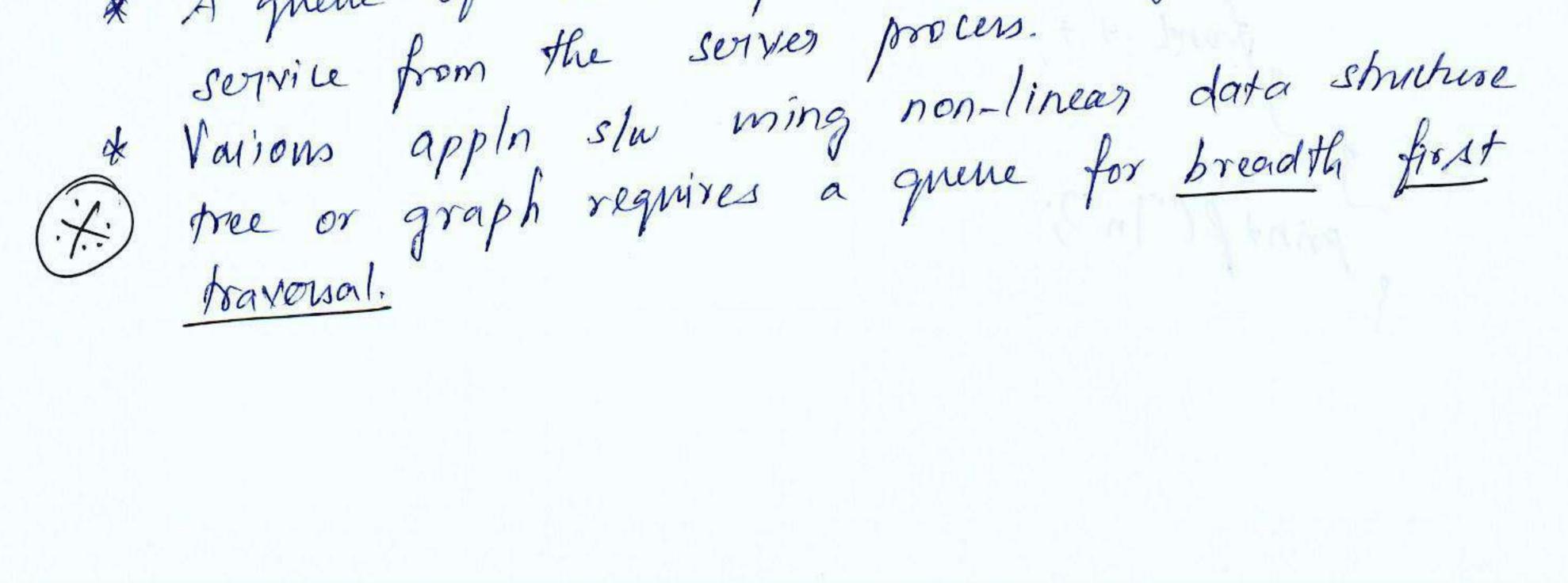


z foont ++; 2 while (front < MAX-1) 2 print & ("1.d", deque (fron £.7); front + +; front = o; while (front 2= rear) printf ("\$i ".d", deque [front]; font + +; prindf("In");

Applications of Queues:

* Queues are widely med as waiting list for a single shared resource like printer, disk, CPU. & " " woed to transfer data asynchronowsky Colata not necessarily received at same rate as sent) b/n two processes (IO buffers), e.g. pipes, file IO, sockets. d" " med as buffers on MP3 players and portable CD players, i Pool playlint. & " " med in Play list for juke box to add songs to the end, play from the front of the list. & " med is operating system for handling interruph. 4 When pomming a real-time system that can be intempted, for example, by a mome click, it is necessary to process the intempts immediately, before proceeding with the current job. If the intempts have to be handled in the order of anival, then a FIFO queue is the appropriate DS. * Scheduling of processes * A queue of client processes waiting to receive the service from the serves process.

65)





Josephus Problem; -

In Jasephun pbm, 'n' people stand in a circle waiting to be executed. * The counting starts at some point in the circle and proceeds in a specific direction (clockwise /anti-clockwise) around the circle. * In each step, a certain no. of people are skipped and the next porron is executed. * The elimination of people makes the circle smaller and smaller. * At the last step, only one person remains who is

de claned the "winner" Thusfore, if there are n'number of people and a number's', which indicates that k. people are skipped and kith person in the circle is eliminated, then the pbm is to choose a position is the initial circle so that the given porson becomes the winner. For example, if there are 5 (or people and every second (k) person is eliminated, then first the person at position 2. In eliminated followed by part position 4 followed by person at position 1 and finally. The person at position 5 is eliminated. ..., the person of position '3' is the winner.

1gm Code:

#include <stdio.h> # include <conio.h> #include <conio.h> #include <malloc.h> somet node ? int player_id; smut node * next; ? Smut node * next; ? Smut node * start, *ptr, * new_node; int main() ? int main() ? int main()? ?

ptrot next = new - node; new - node -7 player - id = i; new - node -7 next = start; ptr = new - node; g for(count = n; count > 1; count -; i < k - 1; + i) i < ptr = ptr -> next;

67)

print f ("In Enter the no. of players:"); Scan f ("y.d", & n); print f ("In Enter the value of k:"); scan f ("y.d", &k); I create circular linked list containing all the players start = malloe (size of (struct node)); Start-7 player_id = 1; ptr=start; for (i=2; i<n; itd) pts-Inext = pts-Inext-Inext; Remove the eliminated player from the circular linked hist y printf("In The winner is player y.d", ptr-splayer.id); return o; y

new-node = malloc (sizeof (smit node));

* - * - * - *